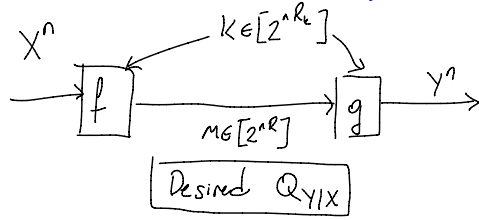


12/13/2016
Tuesday

Distributed Channel Synthesis

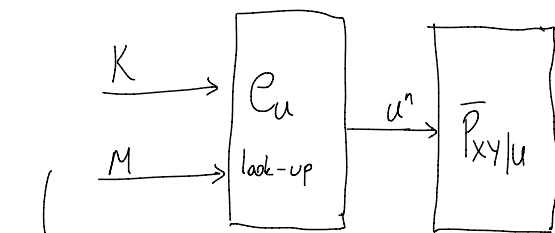


$$\|P_{Y^n|X^n} - P_X Q_{Y|X}\|_{TV} < \epsilon \quad \leftarrow \text{objective}$$

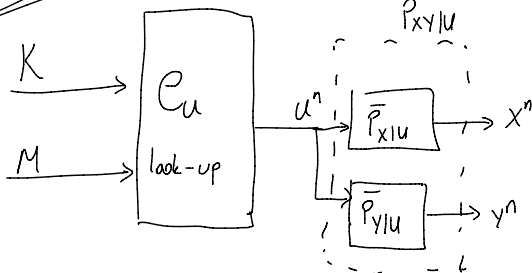
For $R_k = 0$, $R > C_w(X, Y) = \min_{X-U-Y} I(X, Y; U)$
 $\uparrow$
 necessary and sufficient.

$$C_w(X; Y) \in [I(X; Y), \min\{H(X), H(Y)\}]$$

$\uparrow$
let $U=X$ or $U=Y$



ideal distr



$$R_k + R > I(u; X, Y) \quad (\text{soft covering})$$

What's missing: $P_{X^n|K} = P_X$

$\hookrightarrow$

$$R_k + R > I(u; X, Y) \quad (\text{soft covering})$$

$$R > I(u; X)$$

$X^n \perp K \rightarrow$ from block diagram
 $X^n - (M, K) - Y^n$

Recall:

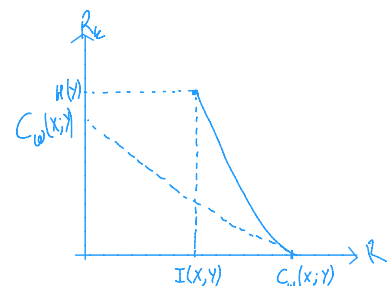
$$C_{GK}(X, Y) = \max_{\substack{U-X-Y \\ X-Y-U}} I(X, Y; U)$$

$$C_{GK}(X; Y) \in [0, I(X; Y)]$$

They look like dual

Note $C_{GK} = I(X; Y) \Leftrightarrow C_w(X; Y) = I(X; Y)$

Choose $\bar{P}_{X,Y|U}$ s.t.
 $X-U-Y$
 and $\bar{P}_{X,Y} = P_X Q_{Y|X}$



Secure DCS

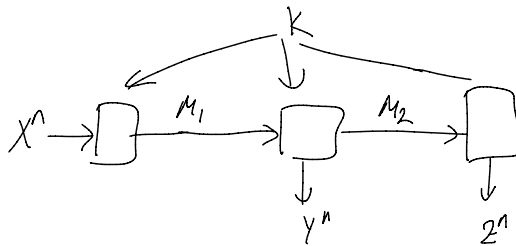
Objective: $\|P_{X^n Y^n U} - P_X^n Q_{Y^n|X} P_U\|_{TV}$ Given U , $X^n Y^n$ look iid and ind. of U

Thm: $R_K > I(U; X, Y)$

for some U : $X-U-Y$

$$R > I(U; X)$$

Aside:



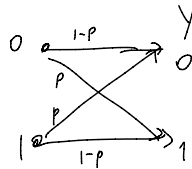
Desired: $Q_{Y^n|X}$

Sai Satpathy solved secure version

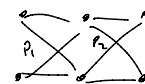
Example

$$X \sim \text{Ber}(1/2)$$

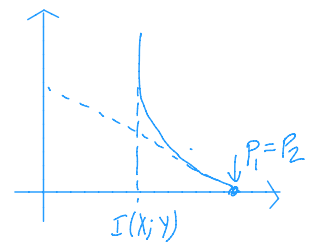
$$Y|X \sim \text{BSC}(p)$$



Optimal U :



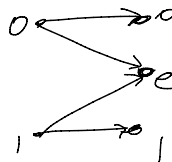
$$p_1 \cdot p_2 = p$$



example:

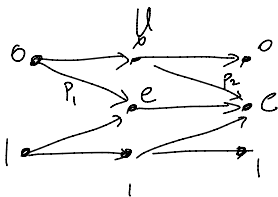
$$X \sim \text{Ber}(1/2)$$

$$Y|X \sim \text{BEC}(p)$$

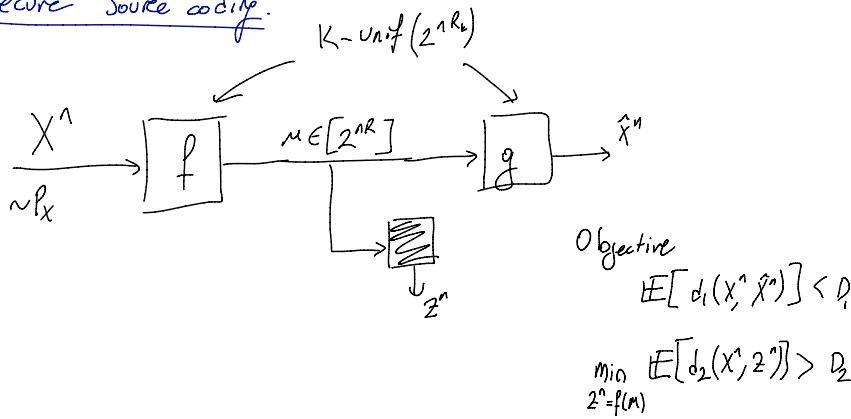


$$p_1 + (1-p_1)p_2 = p$$

$$p_1 + p_2 - p_1 p_2 = p$$



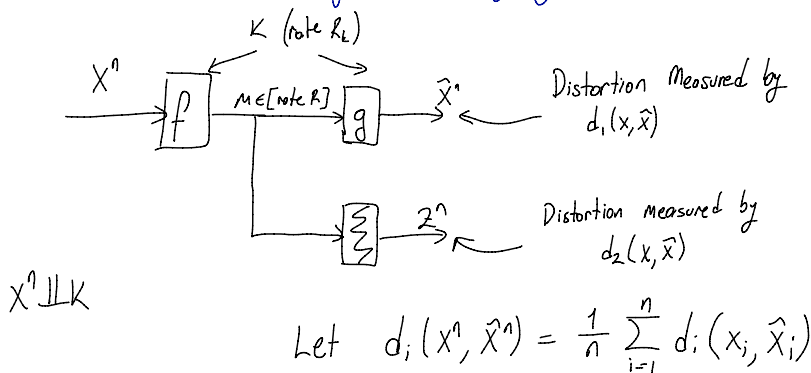
Secure source coding.



$\{(R, R_k, D_1, D_2)\}$ achievable ones.

Rate Distortion Theory for Secrecy Systems:

12/15/2016
Thursday



Performance: $\mathbb{E}[d_1(X^n, \hat{X}^n)] \leq D_1$

$\min_{z(\cdot)} \mathbb{E}[d_2(X^n, z^n)] \geq D_2$

$$\left(= \frac{1}{n} \sum_{i=1}^n \min_{z_i} \mathbb{E}[d_2(X_i, z_i(m))] \right)$$

$R_k = \infty$
because m is indep. of X_i when $R_k = \infty$

Theorem:

Closure of achievable region = $\left\{ (R, R_k, D_1, D_2) : \begin{array}{l} \exists p_{\hat{X}|X} \text{ s.t.} \\ \mathbb{E}[d_1(X, \hat{X})] \leq D_1 \\ \mathbb{E}[d_2(X, \hat{X})] \geq D_2 \\ I(X, \hat{X}) \leq R \end{array} \right\}$

How?

~~$R_k > 0$~~ (already satisfied)
because this is the closure
(we actually need $R_k > 0$)